

# MultiPath TCP : Linux Kernel implementation



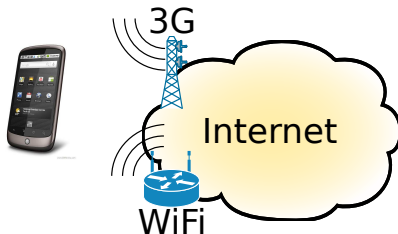
Presenter: Christoph Paasch  
IP Networking Lab  
UCLouvain, Belgium

August 28, 2012

<http://multipath-tcp.org>

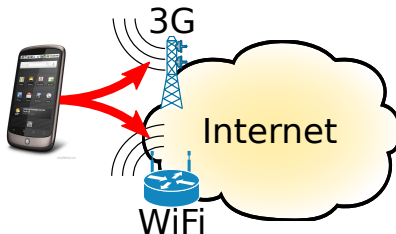
# Networks are becoming Multipath

Mobile devices can connect to the Internet via different interfaces



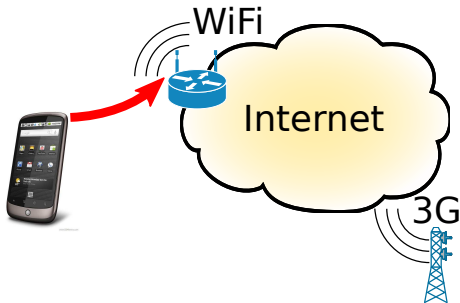
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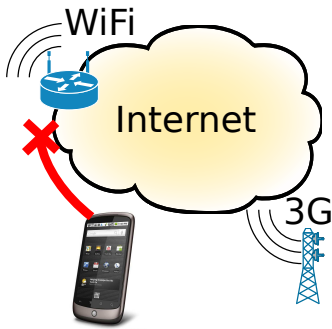
## However, TCP does not support this

Smartphones have to restart their data-transfer when moving away from the WiFi access-point.



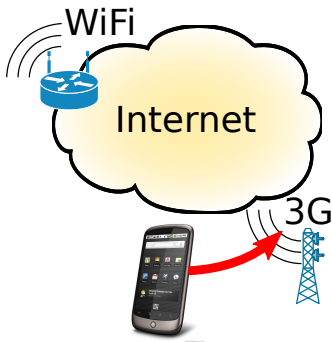
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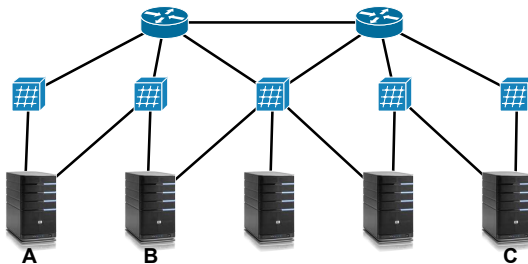
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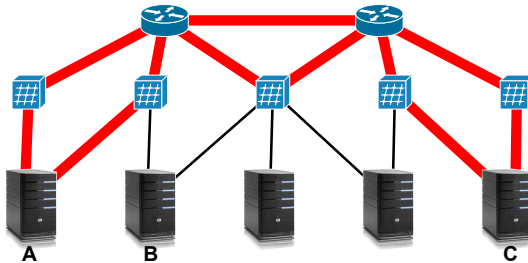
# Networks are becoming Multipath

Data-centers have a large redundant infrastructure



# Networks are becoming Multipath

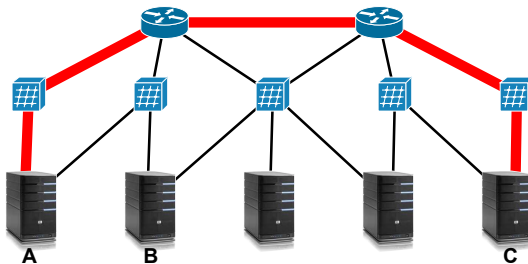
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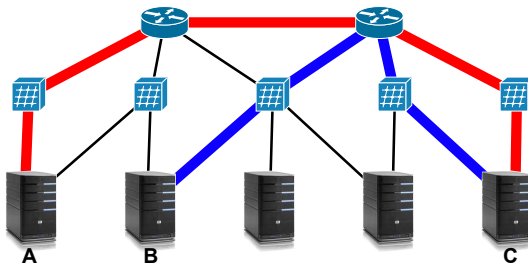
## However, TCP is suboptimal

Collisions in data-center reduce the bandwidth and result in suboptimal load-balancing



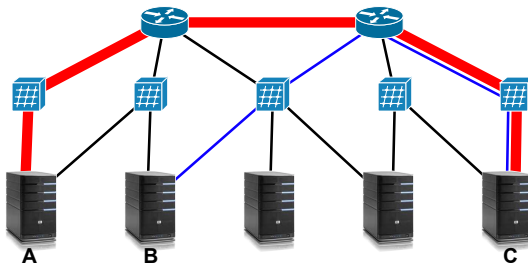
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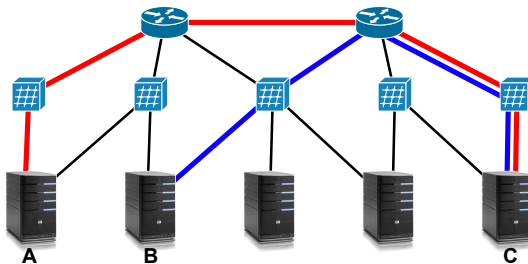
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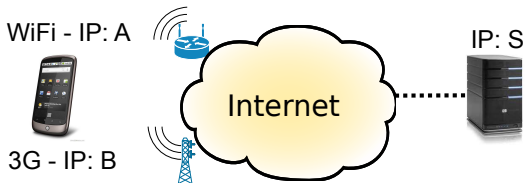


Mismatch between the **multipath network**  
and the **single-path transport** protocol.

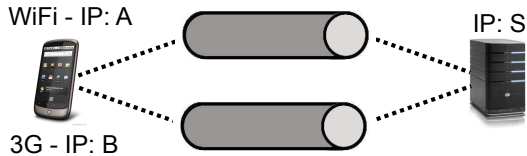
# MultiPath TCP

- Runs with **unmodified applications**
- Works over **today's Internet**
- **IPv4/IPv6** are both supported (even simultaneously)

# MultiPath TCP

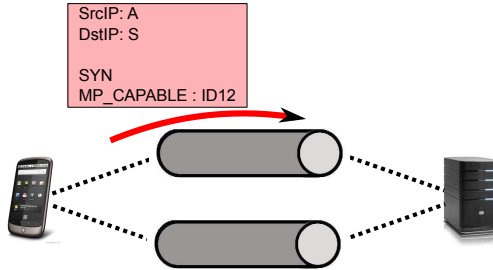


# MultiPath TCP

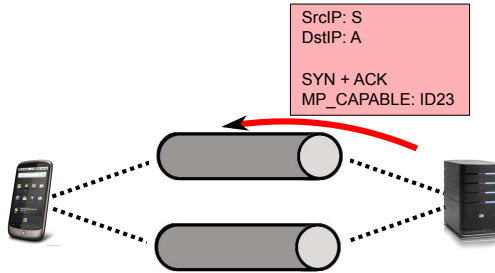




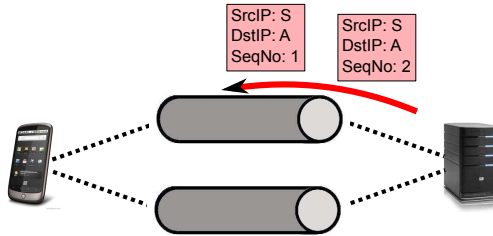
Is the server MPTCP-capable?



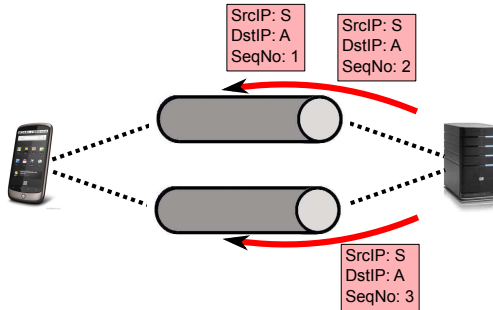
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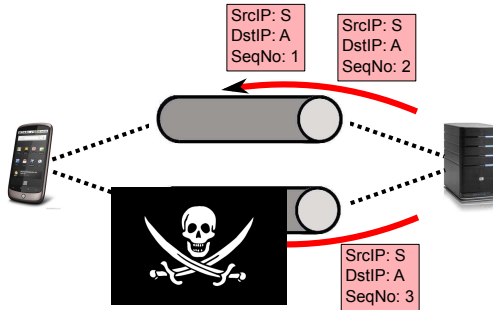
## Sending Data - Naiv approach



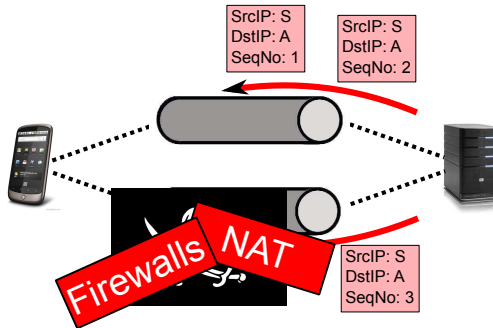
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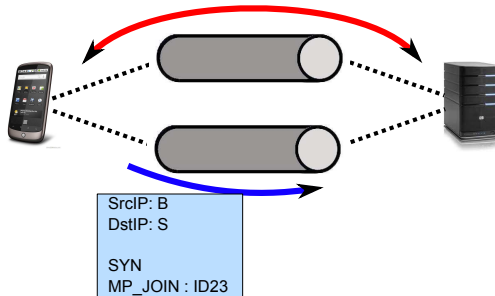
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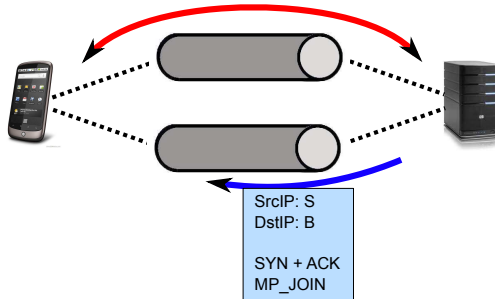
## Sending Data - Naiv approach



## Establish separate subflows

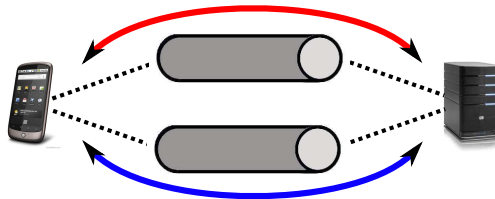


## Establish separate subflows

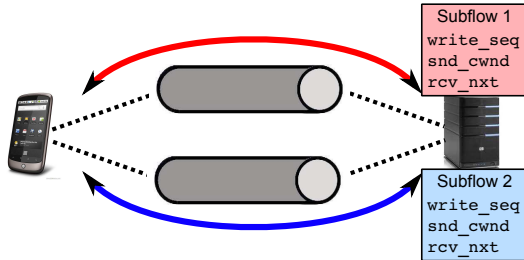




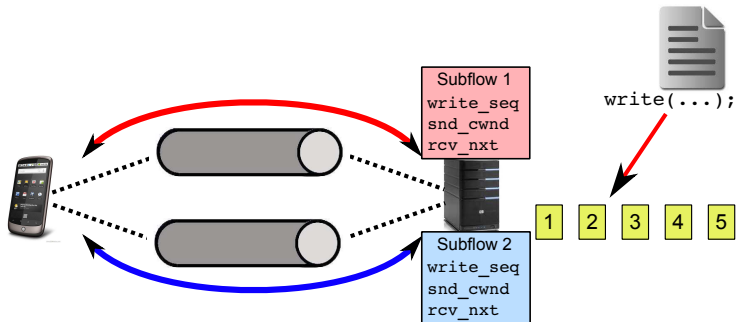
## Establish separate subflows



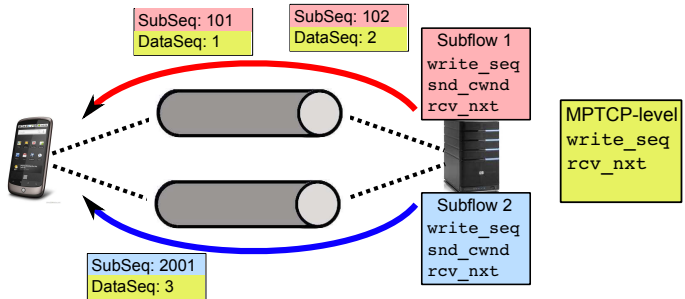
## Establish separate subflows



## Sending Data



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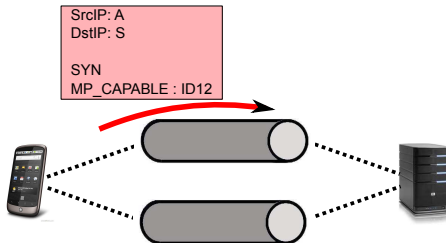
- Both subflows can be used simultaneously.
- Addresses are advertised with ADD\_ADDR and removed by REMOVE\_ADDR.
- Subflows can be dynamically added and removed during the lifetime of the connection.

# MultiPath TCP

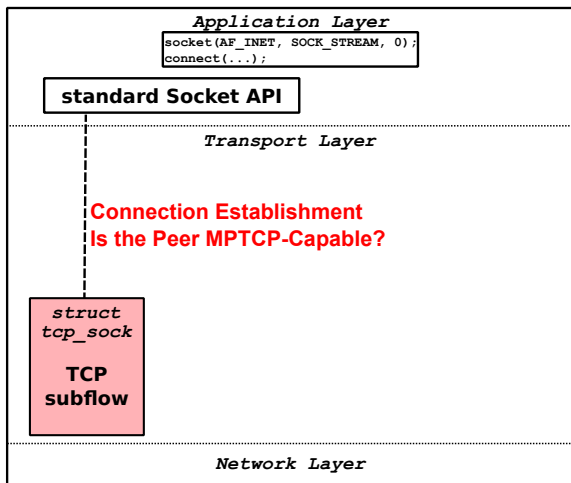
## Linux Kernel Implementation

Available at <http://multipath-tcp.org>

## Exchanged Messages

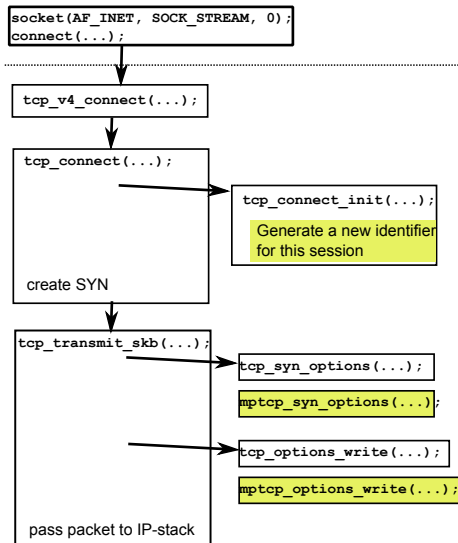


## High-Level Kernel design - Client Side

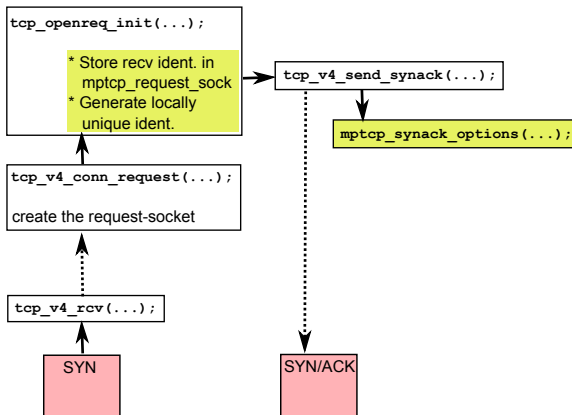




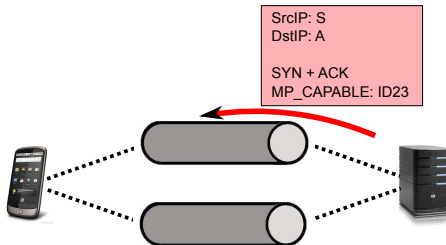
## In-depth call-stack - Client Side



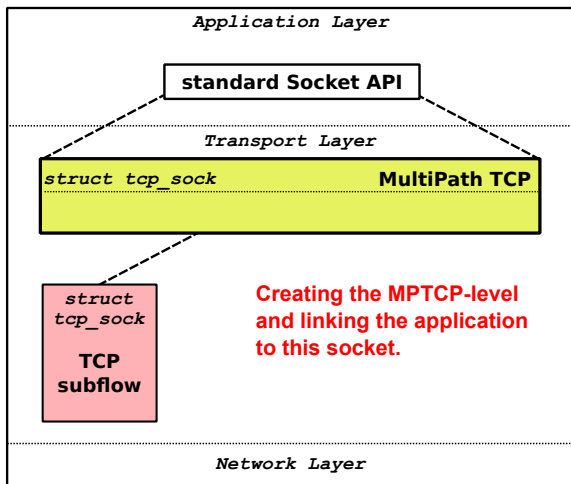
## In-depth call-stack - Server Side



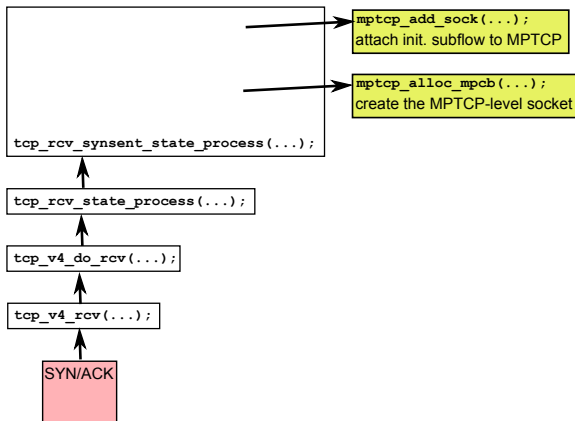
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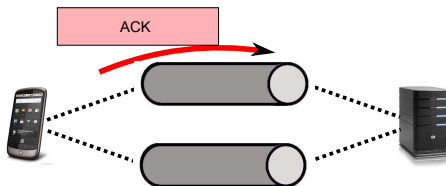
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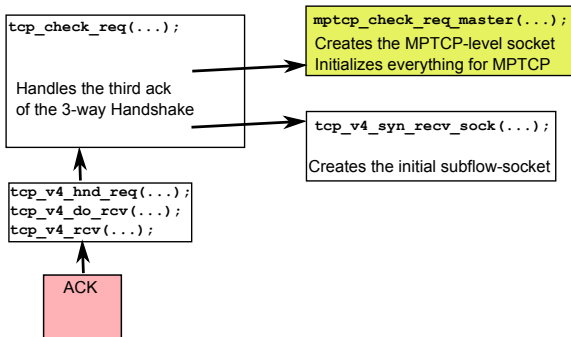
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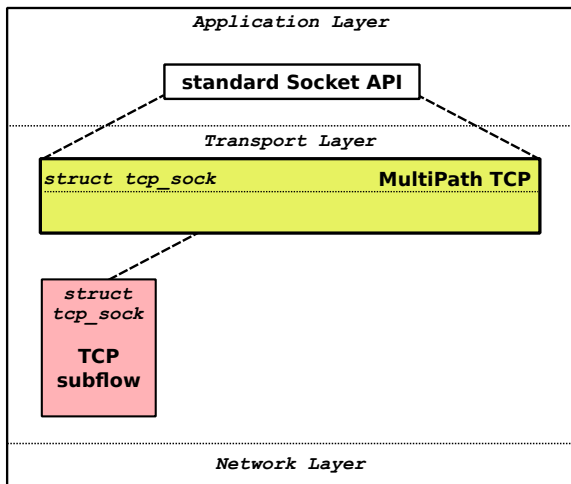
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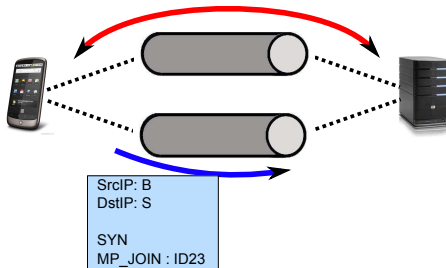


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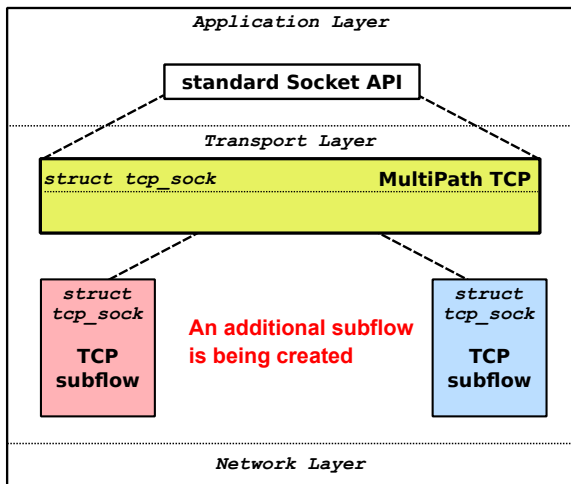




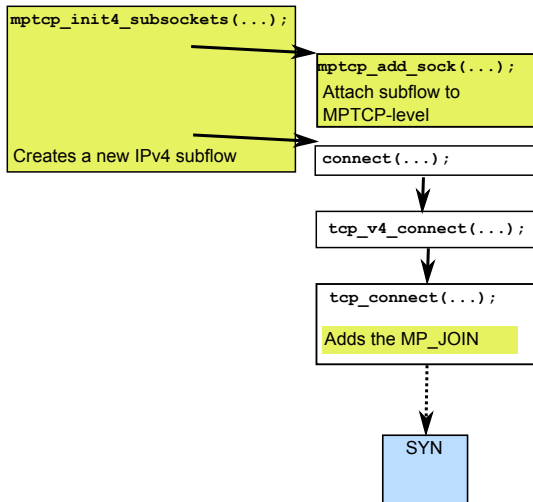
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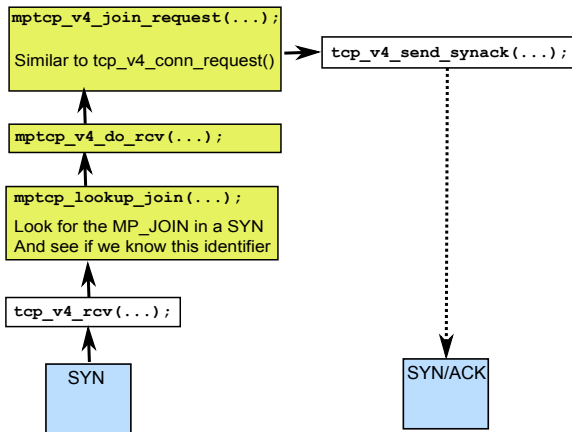
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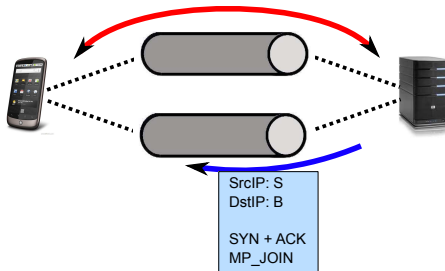
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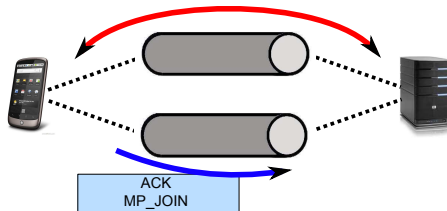
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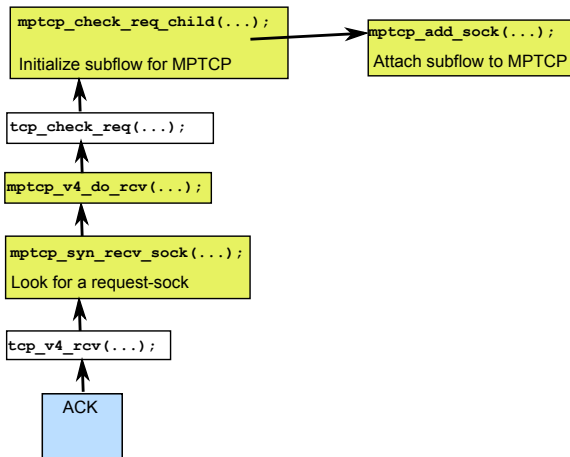
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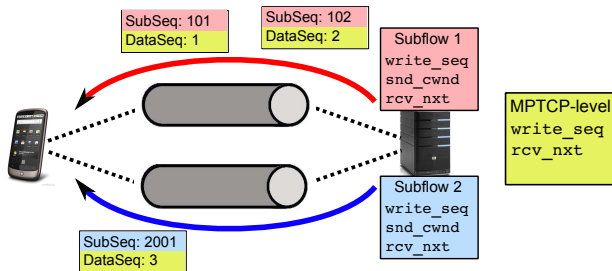
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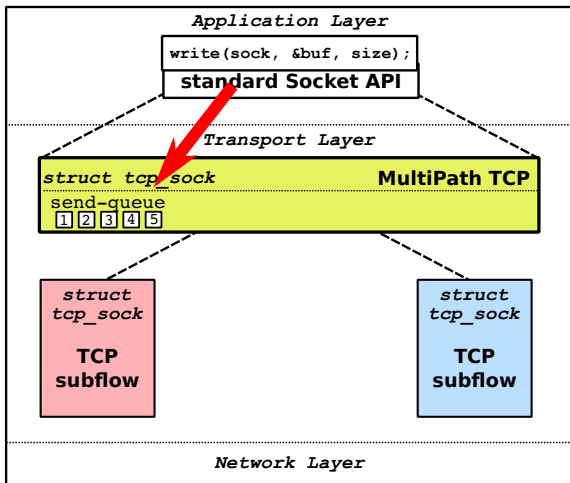


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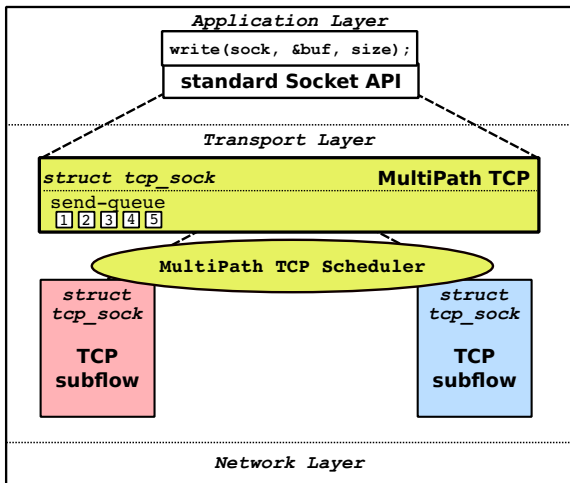




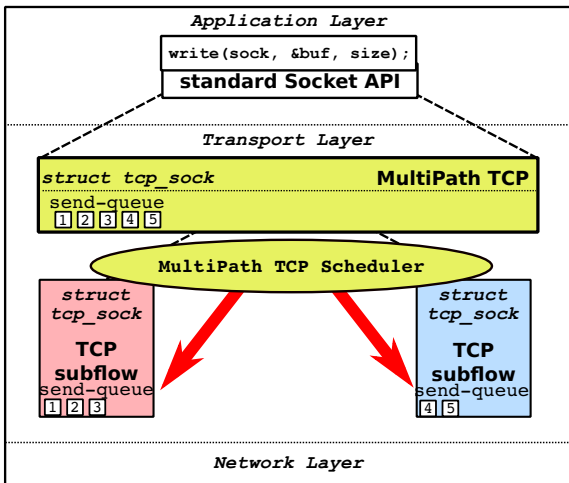
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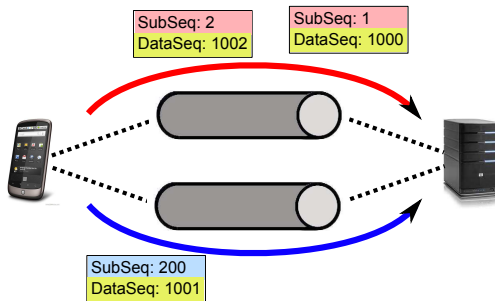


## High-Level Kernel design



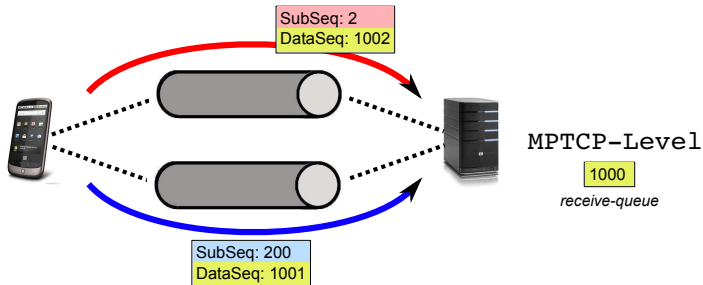
# Receiving Data

Packets can be reordered at the data-level due to delay-differences.



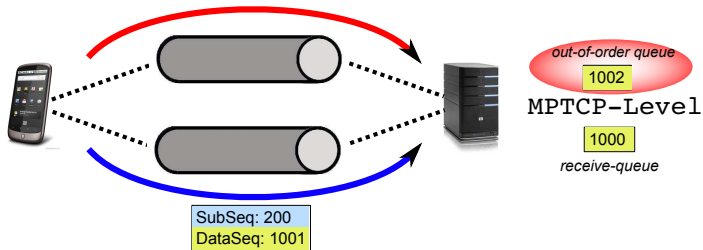
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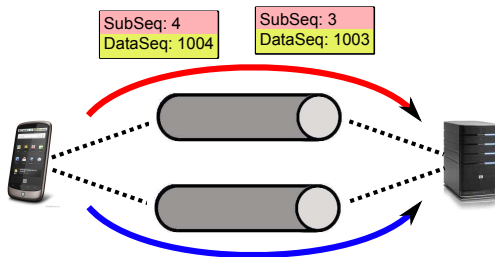
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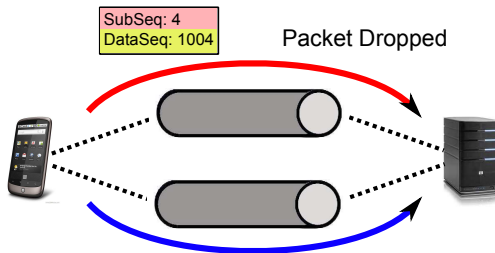
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A loss at the subflow-level (or network-reordering) can also cause reordering at the subflow-level



# Receiving Data

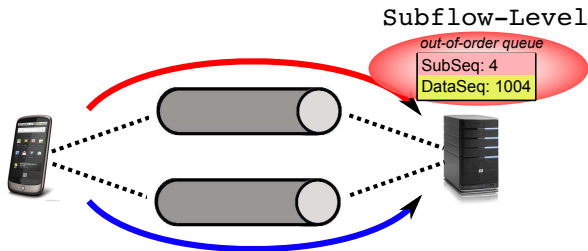
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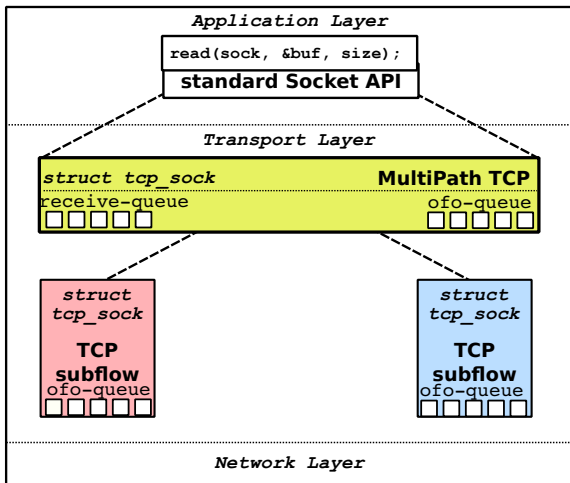


# Receiving Data

Subflow-level out-of-order queues are necessary to handle the retransmission at the subflow-level



## High-Level Kernel design



# MultiPath TCP

## Design Challenges

## *Naive approach*

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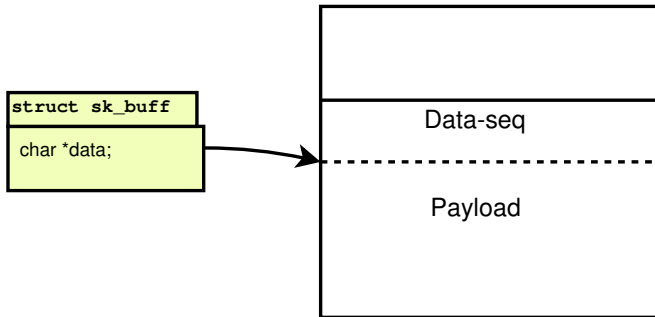
```
struct tcp_skb_cb {  
    __u32          seq; /* Starting sequence number */  
    [...]            
#ifdef CONFIG_MPTCP  
    __u32          data_seq;  
    __u32          end_data_seq;  
    __u32          data_ack;  
    [...]            
#endif  
}
```

---

- Writing data-seq in *tcp\_options\_write*
- Increased *tcp\_skb\_cb* by 24 bytes.

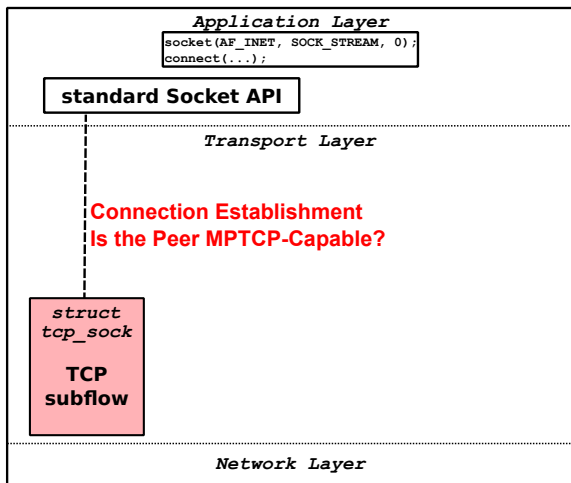
## *Our solution*

Inside the MPTCP-scheduler write the data-seq on top of the payload, **before** calling *tcp\_transmit\_skb*.



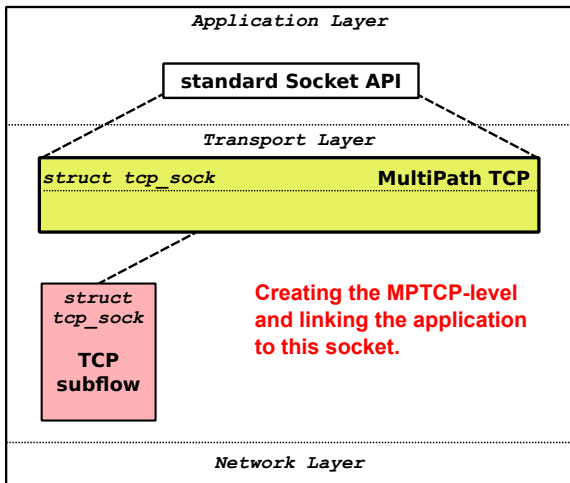
No more increase of *tcp\_skb\_cb*.

## *Current Design*



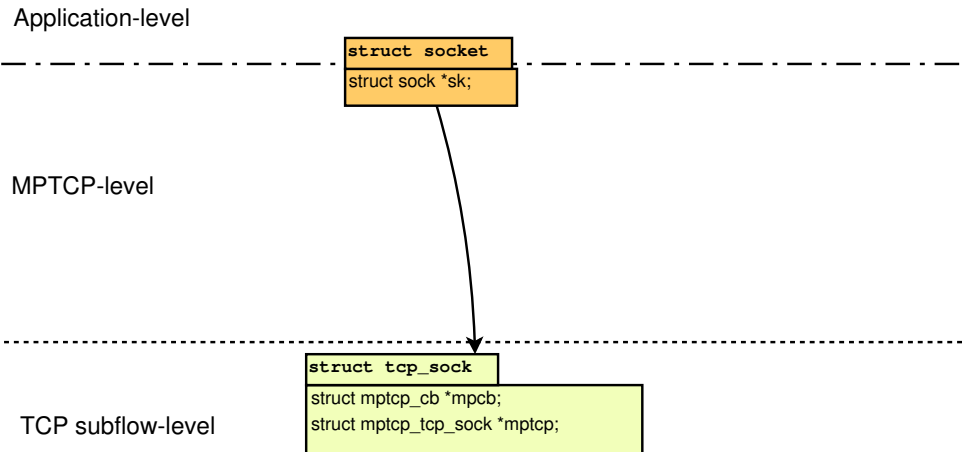
# Creating the MPTCP-level on the client-side.

## *Current Design*



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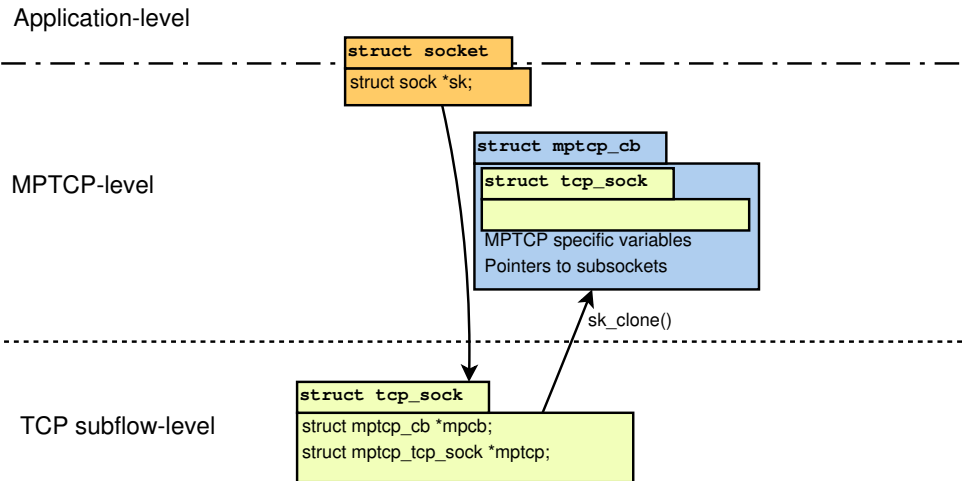
## *Current Design*





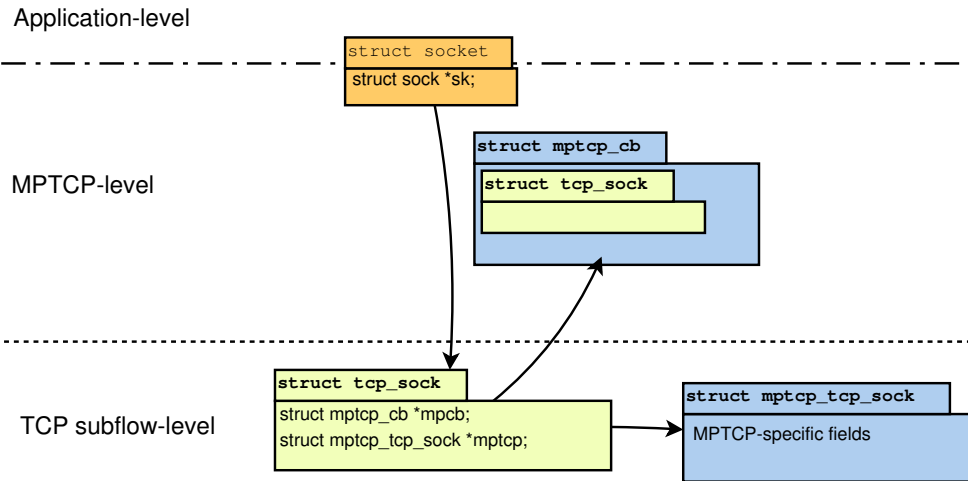
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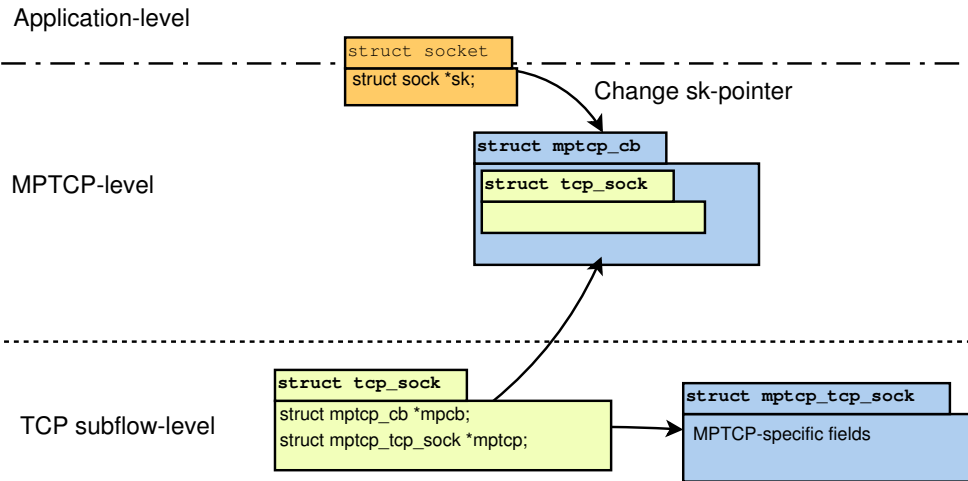
# Creating the MPTCP-level on the client-side.

## *Current Design*



# Creating the MPTCP-level on the client-side.

## *Current Design*



## *Current Design*

- Problems, if the application does a system-call on the socket, **before** the reception of the SYN+ACK
- Fix: Wait for the SYN+ACK. E.g., `tcp_sendmsg`:

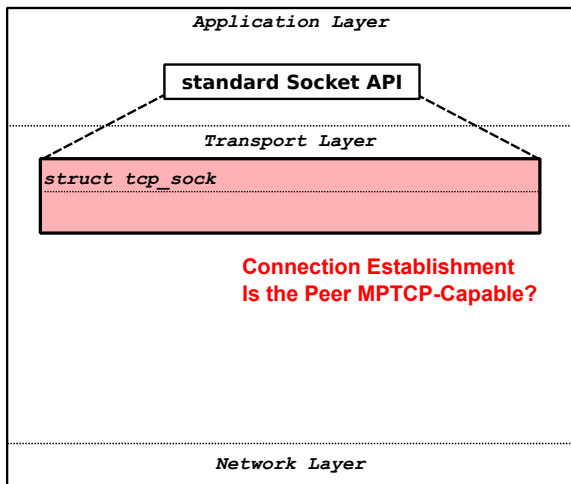
---

```
/* Wait for a connection to finish. */  
if ((1 << sk->sk_state) & ~(TCPF_ESTABLISHED |  
    TCPF_CLOSE_WAIT))  
    if ((err = sk_stream_wait_connect(sk, &timeo))  
        != 0)  
        goto do_error;
```

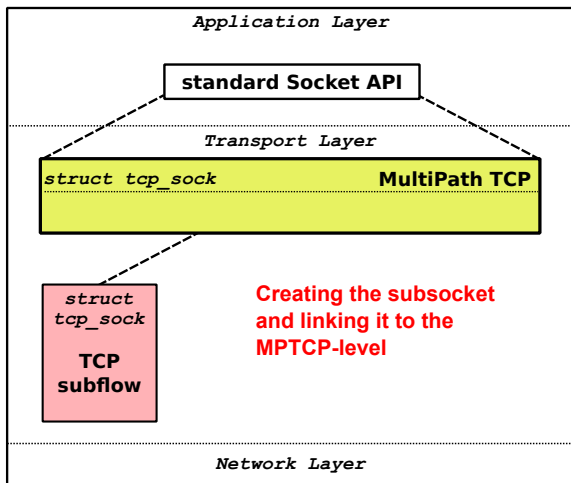
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- We need to do this in **all** functions that take a lock on the socket! `tcp_recvmsg`, `tcp_splice_read`, `ip_setsockopt`, `ip_getsockopt`, `tcp_ioctl`, ...and many more

## Upcoming Design

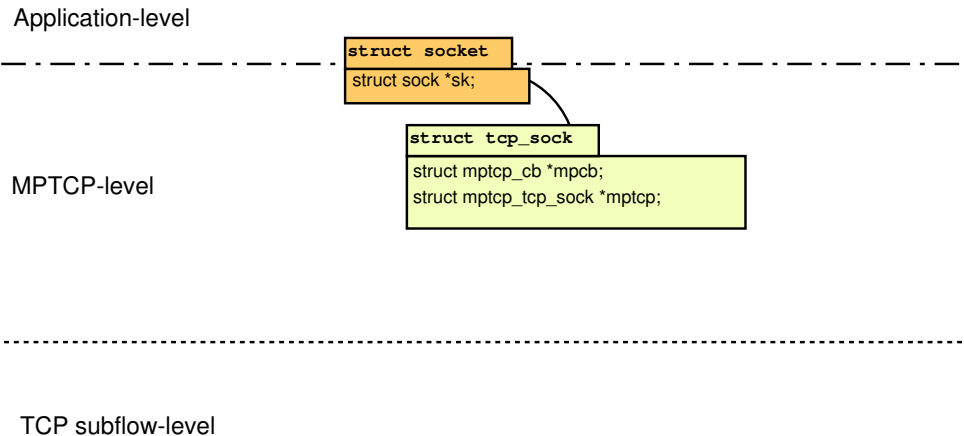


## Upcoming Design



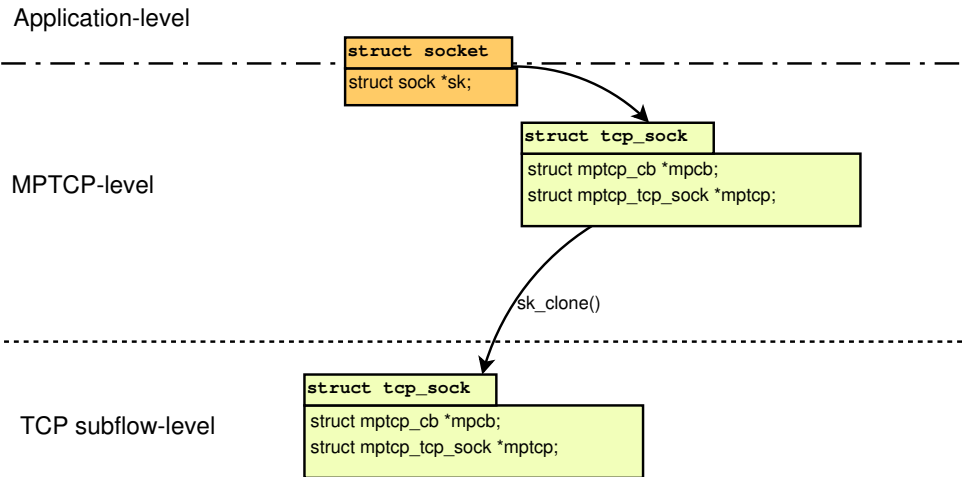
# Creating the MPTCP-level on the client-side.

## *Upcoming Design*



# Creating the MPTCP-level on the client-side.

## *Upcoming Design*





# Creating the MPTCP-level on the client-side.

## *Upcoming Design*

Application-level

```
struct socket
struct sock *sk;
```

MPTCP-level

```
struct mptcp_cb
MPTCP specific variables
Pointers to subsockets
```

```
struct tcp_sock
struct mptcp_cb *mpcb;
struct mptcp_tcp_sock *mptcp;
```

TCP subflow-level

```
struct tcp_sock
struct mptcp_cb *mpcb;
struct mptcp_tcp_sock *mptcp;
```

```
struct mptcp_tcp_sock
Per socket MPTCP-specific
fields
```

## *Questions*

- Lots of socket options in the TCP/IP stack
- Some are for the MPTCP-level (*SO\_SNDBUF*), some should get passed onto all other subflows (*IP\_TTL*)
- This requires a lot of changes in TCP unrelated functions (e.g., *do\_ip\_setsockopt*)

How could we handle this?

## *Questions*

We have a lot of:

---

```
if ( tcp_sk(sk) -> mpc ) {  
    DO_SOME_MPTCP_STUFF  
} else {  
    DO_USUAL_TCP_STUFF  
}
```

---

## Submitting MPTCP upstream???

- ~ 10000 lines of code
- Tightly integrated in the TCP-stack
- More work to do:
  - Cleanup - better separate MPTCP from TCP
  - Some missing features
  - Support TSO
  - Support NET\_DMA
  - ...
- How to split the patch in small pieces?

Freely available at **<http://multipath-tcp.org>**  
Download it, try it out, contribute!

UCLouvain MPTCP-Team:

Sébastien Barré

Christoph Paasch

Gregory Detal

Fabien Duchene

Prof. Olivier Bonaventure

Thanks to our previous and present partners/contributors:

